

**Language Design,
Economics of Programming Languages,
and
Overview of COOL**

Grade Weights

- Project 50%
 - I, II 10% each
 - III, IV 15% each
- Midterm 15%
- Final 25%
- Written Assignments 10%
 - 2.5% each

Lecture Outline

- Today's topic: language design
- Why are there new languages? (The Economy of Languages)
- Good-language criteria
- History of ideas:
 - Abstraction
 - Types
 - Reuse
- Cool
- The Course Project

Programming Language Economics: Three Questions

- Why are there so many programming languages?
 - Hundreds if not thousands in everyday use.
 - Why do we need these?
 - Why not just one all-purpose language?
- Why are there new programming languages given that we already have so many?
- How do we know a good programming language when we see one?

Why So Many Programming Languages?

- » Well, why so many different models of cars?
 - » Fast cars
 - » Cars with good fuel economy
 - » Cars (i.e., trucks, SUVs) that can carry larger amounts of cargo
 - » Cars for various terrains
- » Pretty clear why one model of car would not suffice for all potential uses: No one car can do all things for all drivers.

Programming Domains

- » Scientific Computing (big science, long running apps, simulations)
 - » Good floating point support
 - » Good support for arrays and operations on arrays of floating point data
 - » Good support for parallelism
 - » Many other required qualities
- » Solution: (modern) FORTRAN
 - » still used by many scientists today

Programming Domains

- » Business Computing
 - » Persistence: don't want to lose your data
 - » Good facilities for report generation
 - » Good support for data analysis
- » Solution: SQL and related relational database programming languages

Programming Domains

- » Systems programming (e.g., embedded systems, control systems, operating systems)
 - » Very fine grained control of low level resources
 - » Support for handling time constraints
 - » e.g., network response times
- » Solution: C, C++

Programming Domains

- » Mobile computing
 - » Ease of network functionality
 - » Appropriate mechanisms for inter-process communication
 - » Capability to provide applications for variety of mobile platforms
 - » Differing power and memory capacities, etc.
- » Solution?: Android versions of Java?

Why Are There So Many Programming Languages?

- » Pretty clear that requirements in various domains are completely different
- » Hopefully it's clear that integrating all of the properties into one language (and having that language do a good job with all requirements) is difficult (and likely impossible)

Given all these programming languages,
why the need for new languages?

For a Computer Language to Be Used...

- » Someone has to design it: usually only one or very few people, relatively inexpensive
- » Someone has to build a compiler: usually 10-20 people for a really large compiler, relatively inexpensive

The Real Cost: All the Users and Educating Them

- » Thousands, hundreds of thousands, potentially millions of developers: time and money spent teaching them language is very large cost
- » Cost of classes, textbooks, etc
- » AND since many programmers teach themselves, their time cost is a real economic cost
- » SO, think about cost of educating a million programmers, it's a significant economic investment.

Consequences

- » Even making small changes to an established language used by millions is very expensive
 - » Even small change to interface to compiler is costly in terms of educating users
 - » So, rate of change of language slows.
 - » For most popular languages, this causes ossification

Another Consequence of Cost of Learning Language

- » Paradoxically, it's relatively easy to start a new language
 - » Zero users, so zero training costs at the beginning
 - » Even with few users, costs of teaching them changes in language are not very high
 - » So new language can adapt much more quickly to changing situations (i.e., just not very expensive to experiment with a new language)

So...

- » When a programmer is choosing between an established language that doesn't change quickly, and a new language that has the potential for real gains in productivity in the short term (and at not too much time and expense), then she's going to go with the new language.
- » When? When there is some void that needs to be filled (i.e., current programming languages not meeting some requirement)

Voids?

- » We are in the middle of an information revolution, with new application domains arising all the time.
- » New kinds of programming emerge every few years.
 - » Right now: Mobile applications (and attendant technology to support mobile computing)
 - » A few years ago: the Internet (the phase during which Java came into being)

Summary: Programming Language Economics

- Languages are adopted to fill a void
 - Enable a previously difficult/impossible application
 - Orthogonal to language design quality (almost)
- Observation: Programmer training is the dominant cost of software development
 - Languages with many users are replaced rarely
 - Popular languages become ossified
 - But easy to start in a new niche . . .

One More Consequence

- » New programming languages tend to look a lot like old programming languages
 - » Leverage the knowledge programmers already have about the old language
 - » Classic Example: Java vs. C++
 - » Made entry to Java easy for C++ programmers

Topic: Language Design

- No universally accepted metrics for design
 - Not among programmers
 - Not among researchers in programming language design

Topic: Language Design

- Claim: "A good language is one people use"
 - Not. By this criteria, Visual Basic is the world's best programming language
 - Besides: lot's goes into language adoption (as we've seen) besides technical excellence
 - Language addresses niche -> becomes established in niche -> then inertia sets in
 - Thus Fortran and Cobol are still used, and languages that we could, if starting over today, design much better.

Language Evaluation Criteria

Characteristic	Criteria		
	Readability	Writeability	Reliability
Simplicity	*	*	*
Data types	*	*	*
Syntax design	*	*	*
Abstraction		*	*
Expressivity		*	*
Type checking			*
Exception handling			*

History of Ideas: Abstraction

- Abstraction = detached from concrete details
- Abstraction necessary to build software systems
- Modes of abstraction
 - Via languages/compilers:
 - Higher-level code, few machine dependencies
 - Via subroutines
 - Abstract interface to behavior
 - Via modules
 - Export interfaces; hide implementation
 - Via abstract data types
 - Bundle data with its operations

History of Ideas: Types

- Originally, few types
 - FORTRAN: scalars, arrays
 - LISP: no static type distinctions
- Realization: Types help
 - Allow the programmer to express abstraction
 - Allow the compiler to check against many frequent errors
 - Sometimes to the point that programs are guaranteed "safe"
- More recently
 - Lots of interest in types
 - Experiments with various forms of parameterization
 - Best developed in functional programming

History of Ideas: Reuse

- Reuse = exploit common patterns in software systems
 - Goal: mass-produced software components
 - Reuse is difficult
- Two popular approaches
 - Type parameterization (`List(int)`, `List(double)`)
 - Classes and inheritance: C++ derived classes
 - Combined in C++, Java
- Inheritance allows
 - Specialization of existing abstraction
 - Extension, modification, hiding behavior

Trends

- Language design
 - Many new special-purpose languages
 - Popular languages to stay
- Compilers
 - More needed and more complex
 - Driven by increasing gap between
 - new languages
 - new architectures
 - Venerable and healthy area

Why Study Languages and Compilers ?

5. Increase capacity of expression
4. Improve understanding of program behavior
3. Increase ability to learn new languages
2. Learn to build a large and reliable system
1. See many basic CS concepts at work

Cool Overview

- Classroom Object Oriented Language
- Designed to
 - Be implementable in a short time
 - Give a taste of implementation of modern
 - Abstraction
 - Static typing
 - Reuse (inheritance)
 - Memory management
 - And more ...
- But many things are left out

A Simple Example

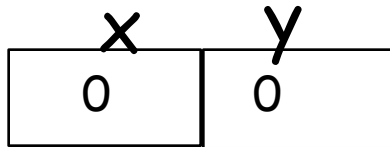
```
class Point {  
    x : Int ← 0;  
    y : Int ← 0;  
};
```

- Cool programs are sets of class definitions
 - A special class **Main** with a special method **main**
 - No separate notion of subroutine
- class = a collection of attributes and methods
- Instances of a class are objects

Cool Objects

```
class Point {  
    x : Int ← 0;  
    y : Int; (* use default value *)  
};
```

- The expression "new Point" creates a new object of class Point
- An object can be thought of as a record with a slot for each attribute



Methods

- A class can also define methods for manipulating the attributes

```
class Point {  
    x : Int ← 0;  
    y : Int ← 0;  
    movePoint(newx : Int, newy : Int): Point {  
        { x ← newx;  
          y ← newy;  
          self;  
        } -- close block expression  
    }; -- close method  
}; -- close class
```

- Methods can refer to the current object using *self*

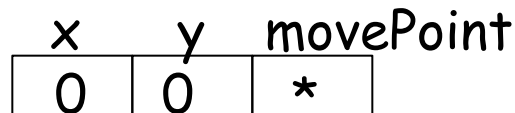
Information Hiding in Cool

- Methods are global
- Attributes are local to a class
 - They can only be accessed by the class's methods
- Example:

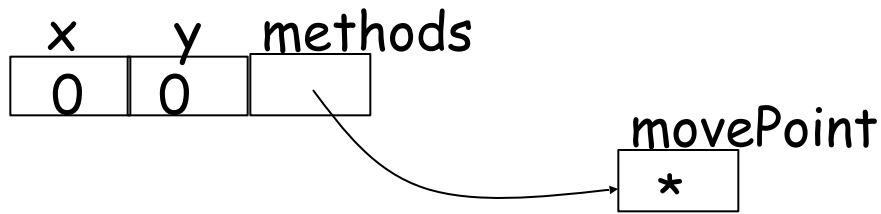
```
class Point {  
    . . .  
    x () : Int { x };  
    setx (newx : Int) : Int { x ← newx };  
};
```

Methods

- Each object knows how to access the code of a method
- As if the object contains a slot pointing to the code



- In reality implementations save space by sharing these pointers among instances of the same class



Inheritance

- We can extend points to colored points using subclassing => class hierarchy

```
class ColorPoint inherits Point {  
  color : Int ← 0;  
  movePoint(newx : Int, newy : Int): Point {  
    { color ← 0;  
      x ← newx; y ← newy;  
      self;  
    }  
  };  
};
```

x	y	color	movePoint
0	0	0	*

Cool Types

- Every class is a type
- Base classes:
 - `Int` for integers
 - `Bool` for boolean values: `true`, `false`
 - `String` for strings
 - `Object` root of the class hierarchy
- All variables must be declared
 - compiler infers types for expressions

Cool Type Checking

```
x : A;  
x ← new B;
```

- Is well typed if **A** is an ancestor of **B** in the class hierarchy
 - Anywhere an **A** is expected a **B** can be used
- Type safety:
 - A well-typed program cannot result in runtime type errors

Method Invocation and Inheritance

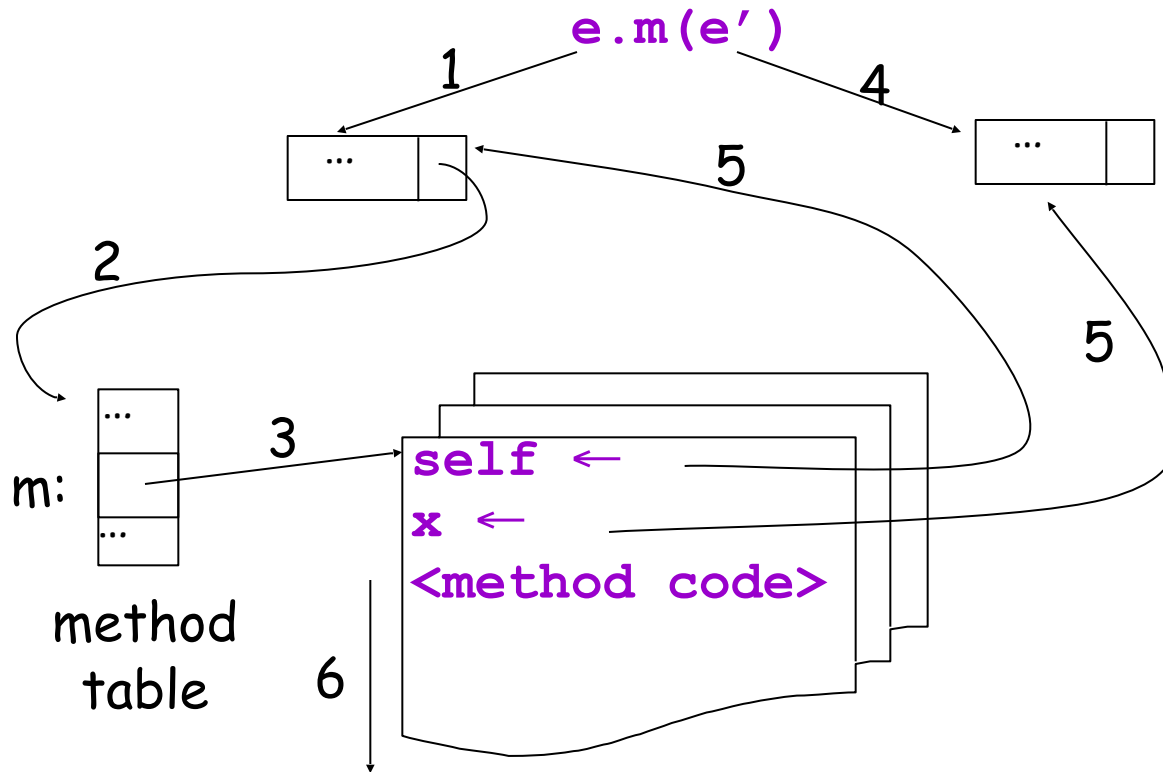
- Methods are invoked by dispatch
- Understanding dispatch in the presence of inheritance is a subtle aspect of OO languages

```
p : Point;  
p ← new ColorPoint;  
p.movePoint(1,2);
```

- `p` has static type `Point`
- `p` has dynamic type `ColorPoint`
- `p.movePoint` must invoke the `ColorPoint` version

Method Invocation

- Example: invoke one-argument method m



1. Eval. e
2. Find class of e
3. Find code of m
4. Eval. argum.
5. Bind self and x
6. Run method

Other Expressions

- Expression language
 - every expression has a type and a value
 - Loops: `while E loop E pool`
 - Conditionals `if E then E else E fi`
 - Case statement `case E of x : Type \Rightarrow E; ... esac`
 - Arithmetic, logical operations
 - Assignment `x \leftarrow E`
 - Primitive I/O `out_string(s), in_string(), ...`
- Missing features:
 - arrays, floating point operations, exceptions, ...

Cool Memory Management

- Memory is allocated every time **new** is invoked
- Memory is deallocated automatically when an object is not reachable anymore
 - Done by the garbage collector (GC)
 - There is a Cool GC

Course Project

- A complete compiler
 - Cool ==> MIPS assembly language
 - No optimizations
- Split in 4 programming assignments (PAs)
- There is adequate time to complete assignments
 - But start early and please follow directions
- Individual or team
 - max. 2 students (but adhere to the Shaw Constraint (SC): can only work with a given partner on one project).